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**An ECC based Anonymous Authentication Protocol for Internet of Things**

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# Outline

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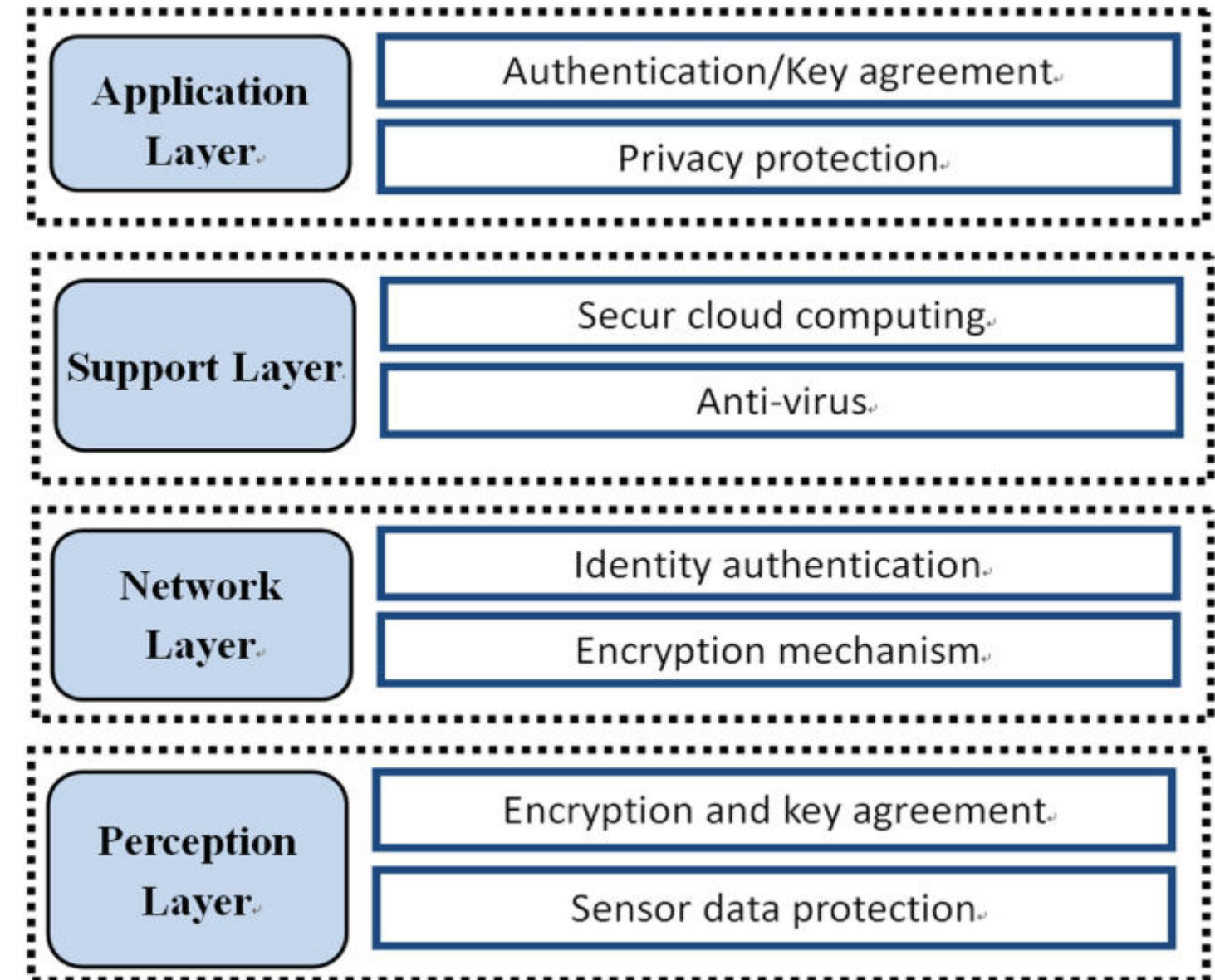
## Introduction

- Advancement of wireless communication and information technologies has resulted in a rapid development of the Internet of Things (IoT).
- In IoT devices all data is transferred through a network between sensing devices, remote users and up to cloud storages.
- The rapid growth in the number of IoTs devices, the heterogeneity and complexity of networks have made **authentication has a challenging task**.
- The simple and low cost nature of IoT devices makes them an attractive target for spoofing or impersonation attacks.
- There are malicious attacks also possible such as impersonation, replay, denial of service, and man-in-the-middle attack.



## Motivation and contribution

- An efficient, secure, and lightweight remote-user authentication-based solution for an IoT environment is necessary.
- Mutual authentication is considered as a key element for successfully accessing various IoT services when it comes to network privacy and security.
- To address this, we propose to design an ECC based anonymous authenticated protocol for internet of things that enables the mutual authentication between users and gateway device.
- The proposed protocol ensures the confidentiality of identity by revealing it exclusively to the server for authentication. No adversary can find the user identity.



IoT Security Architecture

# Elliptic curve cryptography-ECC

- The security of ECC is based on the difficulty of the elliptic curve discrete logarithm problem (ECDLP). The ECDLP is to find  $k$  for given values of  $P$  and  $Q$  in the equation  $Q = kP$ . Finding the value of  $k$  is hard problem.
- ECC based on elliptic curves and requires smaller keys compared to non-elliptic curves cryptography like RSA.
- An elliptic curve  $E_{F_p}$  over a finite field  $F_p$  is defined as the set of all  $(x, y) \in F_p$  such that  $y^2 = x^3 + ax + b$ , where  $a, b \in F_p$  and  $4a^3 + 27b^2 \neq 0$ , along with a distinguished point at infinity which is denoted by  $O$ .

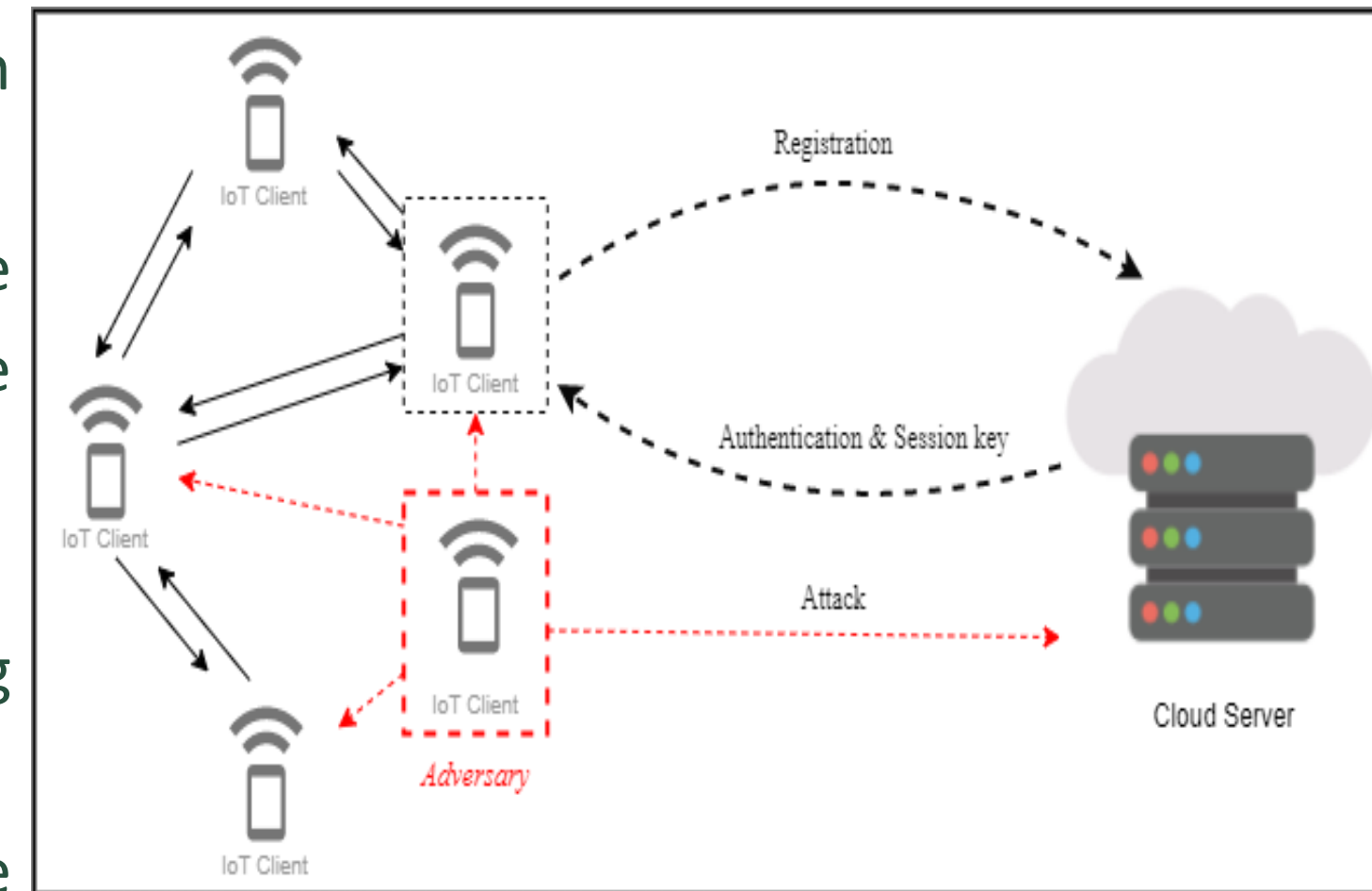
## ***ECC selection criteria***

- *Smaller Key Sizes: ECC keys are much shorter than RSA or DSA keys of equivalent strength.*
- *Faster Encryption and Decryption: ECC algorithms are faster than RSA and DSA algorithms.*
- *Lower Power Consumption: it can help to extend battery life on mobile and IoT devices.*
- *Bandwidth Efficiency: Smaller key sizes also mean that less data needs to be transmitted over the network*

# Proposed Authentication Protocol

## 1. System Model

- The IoT environment scenario consists of two communication parties server S and the IoT device D trying to communicate with authentication in the network.
- An Adversary A with defined capabilities(Dolev-Yaos model) in the IoT network having the access of the communication in the network.
- The server S having enough resources is responsible for allowing genuine device for communication.
- If user wants to communicate with server , it must pass the authentication process.
- The S sends a challenge message to the IoT device , then it authenticated by S and generates a session key once it receives the challenge message. Both S and device D eventually agreed on a shared session key.



IoT network Scenario



# Proposed Authentication protocol Construction

The protocol is carried out in three phases:

## Registration phase:

- At the first step of this phase ,a random nonce  $n_s$  is selected by server(S) as its private key and calculates the curve point  $G_s$  as  $G_s = n_s G$ .
- The server selects the random number  $n_d$  for every device  $D_i$  that request to register on server and verifies with  $R_d = n_d G$ .
- The corresponding values  $R_d, n_d$  for each device  $D_i$  are stored in the database.
- Each device stores  $G_s$  value in the memory

## Algorithm 2 REGISTRATION

Server (S)

Device ( $D_i$ )

- 1: Server generates randomly  $n_s$  (Secret Key)
- 2: Calculate the curve point  $G_s \ni G_s = n_s \times G$
- 3:  $D_i$  randomly generates  $r_i$
- 4: Compute  $ID_i = r_i \times D_i$
- 5: Send  $\langle ID_i \rangle$  to Server S
- 6: Choose Random  $n_d$  for each  $D_i \ni 1 \leq d \leq i$
- 7: Compute  $R_d = n_d \times G; 1 \leq d \leq i$
- 8: Stores  $(n_s, G_s, n_d, R_d,)$
- 9: Send  $\langle G_s, n_d, R_d \rangle$  to each  $D_i$
- 10: Stores  $(G_s, n_d, R_d)$

## Construction-Contd...

### Login phase:

The authentication phase requires Message exchanges and computations between the device and the server as described in the following steps:

- To begin the authentication process, the device  $D_i$  generates a random number  $r_i$  and computes its identity as  $ID_i = (r_i \oplus D_i)$  sends a request to the server as  $\langle \text{Req}; ID_i \rangle$ .

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### Algorithm 3 LOGIN

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Server ( $S$ )	Device ( $D_i$ )
	1: Send $\langle ID_i \rangle$ to Server $S$
2: $S$ retrieves respective $n_d$ & $R_d$	
3: $S$ randomly chooses $n_1$	
4: Compute $R_1 = n_1 \times G$	
5: Compute $R_1' = n_s \times R_1$ $\therefore R_1' = n_1 \times n_s \times G$	
6: Send $\langle R_1' \rangle$ to $D_i$	

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## Algorithm 4 AUTHENTICATION

Server ( $S$ )

Device ( $D_i$ )

- 1:  $D_i$  randomly generates  $n_2$
- 2: Compute  $R_2 = n_2 \times G$
- 3: Compute  $R_3 = n_2 \times R_1'$
- 4: Compute  $R_4 = n_2 \times R_d$
- 5: Compute Authentication Parameter  $V = H(R_3 + R_4)$
- 6: Send  $\langle V, R_2 \rangle$  to Server  $S$
  
- 7: Compute  $V' = H(n_1 \times n_s \times R_2 + n_d \times R_2)$
- 8: **if**  $V == V'$  **then**
- 9:     Authentication Successful
- 10:     $S$  randomly selects  $n_3$
- 11:    Compute  $R_5 = n_3 \times G$
- 12:    Compute  $V_1 = H(n_3 \times R_d)$
- 13:    Compute  $SK = H(R_2 \times n_3 || R_d)$
- 14:    Send  $\langle V_1, R_5 \rangle$  to  $D_i$
- 15: **else**
- 16:     Authentication Failed
- 17: **end if**
  
- 18: Compute  $V_1' = H(n_d \times R_5)$
- 19: **if**  $V_1 == V_1'$  **then**
- 20:     Authentication Successful
- 21:      $SK' = H(R_2 \times n_3 || R_d)$
- 22: **else**
- 23:     Authentication Failed
- 24: **end if**

# Construction-Contd...

## Authentication phase:

# Verification Proof

## Server side Verification

$$V' = H(n_1 \times n_s \times R_2 + n_d \times R_2)$$

$$V' = H(n_1 \times n_s \times n_2 \times G + n_d \times n_2 \times G)$$

$$V' = H(n_2 \times R_1' + n_2 \times R_d) \quad \because R_1' = n_s \times n_1 \times G \quad \because R_d = n_d \times G$$

$$V' = H(R_3 + R_4) = V$$

$$\therefore V' = V$$

## Device side Verification

$$V_1' = H(n_d \times R_5)$$

$$V_1' = H(n_d \times n_3 \times G) \quad \because R_5 = n_3 \times G$$

$$V_1' = H(n_3 \times R_d) = V_1 \quad \because R_d = n_d \times G$$

$$\therefore V_1' = V_1$$

## Session Key Verification

$$SK' = H(R_2 \times n_3 \parallel R_d) = SK$$

$$\therefore SK' = SK$$

## Security and Comparative analysis

- The security level of existing ECC-based authentication protocols are analyzed and Understood that most of the existing protocols are vulnerable to the trace ability attack.
- The robustness of the proposed protocol is evaluated based on formal analysis method using widely accepted approaches based on AVISPA tool.

Scheme	Method	P1	P2	P3	P4	P5
Liao and Hsiao [2]	ECC	X	✓	X	X	X
KS [4]	ECC+Hash	X	X	X	X	X
CWS [7]	ECC+Hash	✓	✓	X	✓	X
KKD [32]	ECC+Hash	X	✓	X	✓	X
WCF [25]	ECC+Hash	X	✓	✓	✓	X
T.M.Butt[23]	ECC+Hash	✓	X	X	X	X
<b>Our method</b>	<b>ECC+Hash</b>	✓	✓	✓	✓	✓

P1-MIM attack, P2-Replay attack, P3-Impersonation attack,  
P4-Message Integrity attack, P5-Traceability attack.  
✓ - Resistant, X - Non-resistant.



## Conclusion and Future work

- *The ECC based authentication is considered for protocol due to its less memory requirements and computational power, eventually suitable for IoT environment.*
- *We proposed a authentication protocol and the security analysis demonstrates that the proposed protocol is provably secure and meets the requirements for security in an IoT environment.*
- *Our future plan is work on lightweight authenticated key agreement protocol for IoT environments.*

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# THANK YOU